Distributed Maintenance of Anytime Available Spanning Trees in Dynamic Networks

A. Casteigts, S. Chaumette, F. Guinand, Y. Pigné

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Distributed Computing



Collaboration of distinct entities to perform a common task.

No centralization available. Direct interaction only.

(Think globally, act locally)

Leader election



Distinguishing exactly one node among all.





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Spanning tree

Selecting a cycle-free set of edges that interconnects all nodes.







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Broadcast

Propagating a piece of information from one node to all others.





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Propagating a piece of information from one node to all others.

Counting



Determining how many participants there are.





Counting



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Consensus, naming, routing, exploration, ...

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In fact, *highly* dynamic networks.



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How changes are perceived ?

- Faults and Failures ?-
- Nature of the system. Change is normal.
- Network is partitioned most of the time.



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Decision should be purely local!

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Coarse-grain model

 \rightarrow Pairwise atomic interaction

(Graph relabeling systems *(Litovsky et al., 1999)*; Population protocols *(Angluin et al., 2004)*)

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Scope of the models

Relations between them (Chalopin, 2006)

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initial states :

I for every node,

- T : a token is on this node
- N : no token is on this node
- \bullet \rightarrow : relation from child to parent



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Properties that hold permanently :

• Each node belongs to exactly one tree



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How about performance?

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Metric of interest?

1. Convergence rate (though not expected to converge)

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... performance evaluation

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 Analysis (coalescing particles in evolving graphs)

... performance evaluation

- Analysis (coalescing particles in evolving graphs)
- Simulations (some are in the 2009 technical report, much more needed)

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... algorithmic aspects

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. performance evaluation

- Analysis (coalescing particles in evolving graphs)
- Simulations (some are in the 2009 technical report, much more needed)

... algorithmic aspects

- Finer-grain adaptation of the principle
 - Synchronous message passing (OK)
 - Semi-synchronous message passing (under study)

. performance evaluation

- Analysis (coalescing particles in evolving graphs)
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... algorithmic aspects

- Finer-grain adaptation of the principle
 - Synchronous message passing (OK)
 - Semi-synchronous message passing (under study)
- Optimization strategies (e.g. from simple Tabou search to full Propp machines)

Thank you !



References :



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